THE QUANTUM METROPOLIS ALGORITHM

An implementation of Metropolis' famous algorithm on a quantum computer.

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Exact calculation of expectation values in thermal equilibrium

$$\langle O \rangle = \frac{\int \mathrm{d}\Gamma \ O(\Gamma) \, \mathrm{e}^{-\beta \, H(\Gamma)}}{\int \mathrm{d}\Gamma \, \mathrm{e}^{-\beta \, H(\Gamma)}} \qquad \xrightarrow{\text{discretization}} \qquad \frac{\sum_{\mu} O(\mu) \, \mathrm{e}^{-\beta \, H(\mu)}}{\sum_{\mu} \mathrm{e}^{-\beta \, H(\mu)}}$$

Problem: Exact evaluation is numerically so absurdly laborious, that it is applicable only for the smallest systems!

Example: A system of 100 two-state particles has 2^{100} different states. The fastest computers do about 10^{16} simple operations per second.

$$\frac{2^{100}}{10^{16} \text{s}^{-1}} = 1.27 \cdot 10^{14} \text{s} = 4 \text{ million years}$$

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Simple Sampling Monte Carlo

$$\langle O \rangle \approx \langle O \rangle_{ss} = \frac{\sum_{\mu_{i=1}}^{\mu_{N}} O(\mu_{i}) e^{-\beta H(\mu_{i})}}{\sum_{\mu_{i=1}}^{\mu_{N}} e^{-\beta H(\mu_{i})}}$$

Problem: A lot of computer time is wasted sampling states that don't add significantly to the sum.

Example: Think of a system at very low temperature, where only a few states near the ground-state contribute to the sums and all the others are exponentially surpressed.

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"Instead of choosing configurations randomly, then weighting them with $\exp(-E/kT)$, we choose configurations with a probability $\exp(-E/kT)$ and weight them evenly."

— Metropolis et al.

Importance Sampling Monte Carlo

$$\langle O \rangle \approx \langle O \rangle_{is} = \frac{1}{N} \sum_{\mu_{i=1}^{(\rho)}}^{\mu_N^{(\rho)}} O\left(\mu_i^{(\rho)}\right) \quad \text{with} \quad \rho = \frac{1}{Z} e^{-\beta H\left(\mu_i^{(\rho)}\right)}$$

Note: This is exactly like an experimental measurement, with the exception that <u>we</u> have to generate the states ourselves, while <u>nature</u> does it for the experimental physicist.

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... a very quick derivation

Problem

It it is not trivial to generate states that follow the Boltzmann distribution!

Idea

- I Find one of the "important states".
- 2 Modify it a little bit to get another important state

Goal

Build a Markov process whose Markov chain contains Boltzmann distributed states.

The Classical Metropolis Algorithm ... a very quick derivation

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... a very quick derivation

Master Equation

Look at the probability $P_k(\mu)$ of the system to be in the state μ after k Markov steps.

$$P_{k+1}(\mu) = P_k(\mu) + \sum_{\nu} \left(P_k(\nu) \ T(\nu \to \mu) - P_k(\mu) \ T(\mu \to \nu) \right)$$

Detailed Balance

In the end, we want $P_k(\mu)$ to be the states Boltzmann probability for all steps k. This implies $P_{k+1}(\mu) = P_k(\mu)$

$$\Rightarrow$$
 $P_k(\nu) \ T(\nu \to \mu) = P_k(\mu) \ T(\mu \to \nu)$

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Transition probabilities

Detailed balance gives the required transition probabilities.

$$\frac{T(\mu \to \nu)}{T(\nu \to \mu)} = \frac{P(\nu)}{P(\mu)} = \frac{Z^{-1} e^{-\beta H(\nu)}}{Z^{-1} e^{-\beta H(\mu)}} = e^{-\beta (H(\nu) - H(\mu))}$$

Two step Markov process

$$\frac{T(\mu \to \nu)}{T(\nu \to \mu)} = \frac{S(\mu \to \nu)}{S(\nu \to \mu)} \frac{A(\mu \to \nu)}{A(\nu \to \mu)} = e^{-\beta (H(\nu) - H(\mu))}$$

Where $S(\mu \to \nu)$ is the probability that a transition is suggested and $A(\mu \to \nu)$ is its probability to be accepted.

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The Classical Metropolis Algorithm ... a very quick derivation

Choice of $S(\mu \to \nu)$ and $A(\mu \to \nu)$

Choose $S(\mu \to \nu) = S(\nu \to \mu)$ equal in both directions.

$$\Rightarrow \frac{A(\mu \to \nu)}{A(\nu \to \mu)} = e^{-\beta (H(\nu) - H(\mu))}$$

Choose the largest acceptance probabilities that still satisfy the equation.

$$A(\mu \to \nu) = \min \left(1, e^{-\beta \left(H(\nu) - H(\mu)\right)}\right)$$

This particular choice of S and A defines the Metropolis Algorithm!

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... an example: 2d Ising-Model

The Ising-Model Hamiltonian

$$H = -J\sum_{\langle ij\rangle} S_i S_j - B\sum_i S_i$$

Implementation of the Metropolis Algorithm

- 1 Start in an arbitrary state.
- 2 E.g.: Randomly select a single spin to flip.
- \blacksquare Calculate the resulting energy difference ΔE .
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The Basics of Quantum Computing ... bits and qubits

Classical Bit

- bit = binary digit or basic indissoluble information unit
- basic unit for measuring information
- has the two states 0 and 1

Quantum Bit

- \blacksquare qubit = quantum bit
- basic unit of quantum information
- state in a two dimensional Hilbert space

$$|\psi\rangle = c_0 |0\rangle + c_1 |1\rangle$$
 with $|c_0|^2 + |c_1|^2 = 1$

 $| 0 \rangle$ and $| 1 \rangle$ are called the computational basis

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The Basics of Quantum Computing

... registers and quantum registers

Registers

- \blacksquare a collection of N classical bits, e.g. (01101010)
- \blacksquare has 2^N different states

Quantum Registers

- a tensor product of N qubits, e.g. $|\phi\rangle = |\psi\rangle_1 \otimes |\psi\rangle_2$
- \blacksquare state in a 2^N dimensional Hilbert space, e.g.

$$|\phi\rangle = c_{00} |00\rangle + c_{01} |01\rangle + c_{10} |10\rangle + c_{11} |11\rangle$$

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The Quantum Metropolis Algorithm ... what is it?

Goal

We want a quantum computer to sample the energy eigenstates $|\psi_i\rangle$ of a given Hamiltonian according to the Boltzmann distribution.

Straightforward translation into quantum mechanics

- I Start in a random energy eigenstate $|\psi_i\rangle$ with energy E_i .
- 2 Suggest a nearby energy eigenstate $|\psi_j\rangle$ with energy E_j .
- 3 Calculate their energy difference $\Delta E = E_j E_i$
- 4 Go to $|\psi_j\rangle$ with probability min $(1, e^{-\beta \Delta E})$.
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... unfortunately there are some problems with this ...

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Quantum Phase Estimation Algorithm

Attaches the binary representation of the states energy as a quantum register.

$$\sum_{i} c_{i} |\psi_{i}\rangle \longmapsto \sum_{i} c_{i} |\psi_{i}\rangle \otimes |E_{i}\rangle$$

Reading of the energy register collapses the state to the corresponding energy eigenstate.

- 1 Preparing a random energy eigenstate.
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There is no way to go from one eigenstate to the other directly!

Generation of new eigenstates

Apply a random <u>local</u> unitary transformation C.

$$C: |\psi_i\rangle \longmapsto \sum_j c_j |\psi_j\rangle$$
 where $E_j \approx E_i$

Next step? Use QPE to collapse to a new eigenstate $|\psi_j\rangle$ and learn its energy E_j ?

$$\sum_{j} c_{j} |\psi_{j}\rangle \longmapsto \sum_{j} c_{j} |\psi_{j}\rangle \otimes |E_{j}\rangle \longmapsto |\psi_{j}\rangle \otimes |E_{j}\rangle$$

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Measuring E_j irreversibly collapses into state $|\psi_j\rangle$! How would we go back to $|\psi_i\rangle$ if we reject the move?

A less destructive measurement

Idea: An energy measurement reveals more information than we actually need! One bit (accept/reject) would be enough ...

$$\sum_{j} c_{j} |\psi_{j}\rangle \otimes |E_{j}\rangle \longmapsto \sum_{j} c_{j} |\psi_{j}\rangle \otimes |E_{j}\rangle \otimes \left(\sqrt{w_{ij}} |1\rangle + \sqrt{1 - w_{ij}} |0\rangle\right)$$

If we measure the last qubit, only one bit of information is revealed and less damage is done to the state.

Better chance to undo it!

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Two different binary measurements

We can use QPE to determine if we are back in state $|\psi_i\rangle$. Formally, we can define a projector on $|\psi_i\rangle$.

$$P = 1 |\psi_i\rangle \langle \psi_i| + 0 |\psi_i^{\perp}\rangle \langle \psi_i^{\perp}| = |\psi_i\rangle \langle \psi_i|$$

Use accept/reject as another binary measurement

$$Q^{\parallel} = 1 |1\rangle \langle 1| + 0 |0\rangle \langle 0| = |1\rangle \langle 1|$$

$$Q^{\perp} = 0 |1\rangle \langle 1| + 1 |0\rangle \langle 0| = |0\rangle \langle 0|$$

$$Q = Q^{\parallel} + Q^{\perp} = \mathbb{1}$$

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$$P = 1 |\psi_i\rangle \langle \psi_i| + 0 |\psi_i^{\perp}\rangle \langle \psi_i^{\perp}| = |\psi_i\rangle \langle \psi_i|$$

Use accept/reject as another binary measurement.

$$\begin{split} Q^{\parallel} &= 1 \left| 1 \right\rangle \left\langle 1 \right| + 0 \left| 0 \right\rangle \left\langle 0 \right| = \left| 1 \right\rangle \left\langle 1 \right| \\ Q^{\perp} &= 0 \left| 1 \right\rangle \left\langle 1 \right| + 1 \left| 0 \right\rangle \left\langle 0 \right| = \left| 0 \right\rangle \left\langle 0 \right| \\ \Rightarrow \quad Q &= Q^{\parallel} + Q^{\perp} = \mathbb{1} \end{split}$$

Damage to the system's state

A measurement of the accept/reject qubit also has an effect on the first qubit register, that holds the system's state.

$$\sum_{j} c_{j} |\psi_{j}\rangle \xrightarrow{Q^{\parallel}} \begin{cases} |\phi_{Q}^{\parallel}\rangle & \text{if accepted} \\ |\phi_{Q}^{\perp}\rangle & \text{if rejected} \end{cases}$$

We can express the initial state $|\psi_i\rangle$ as a superposition of the states $|\phi_Q^{\perp}\rangle$ and $|\phi_Q^{\parallel}\rangle$.

$$|\psi_i\rangle = \left(Q^{\parallel} + Q^{\perp}\right)|\psi_i\rangle = \sqrt{q}|\phi_Q^{\parallel}\rangle + \sqrt{1-q}|\phi_Q^{\perp}\rangle$$

... the implementation

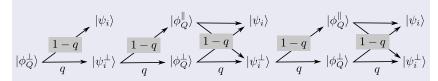
Q^{\parallel} measurement

$$|\psi_i\rangle = \sqrt{q}|\phi_Q^{\parallel}\rangle + \sqrt{1-q}|\phi_Q^{\perp}\rangle$$
$$|\psi_i^{\perp}\rangle = \sqrt{1-q}|\phi_Q^{\parallel}\rangle - \sqrt{q}|\phi_Q^{\perp}\rangle$$

P measurement

$$\begin{aligned} |\phi_Q^{\parallel}\rangle &= \sqrt{q} \, |\psi_i\rangle + \sqrt{1-q} |\psi_i^{\perp}\rangle \\ |\phi_Q^{\perp}\rangle &= \sqrt{1-q} \, |\psi_i\rangle - \sqrt{q} |\psi_i^{\perp}\rangle \end{aligned}$$

Iteration of $P[Q^{\parallel}]$ measurements



Probability never to hit $|\psi_i\rangle$ goes down exponentially with the number of iterations.

Quantum Metropolis Algorithm (1/3)

Measure the energy of an arbitrary initial state using QPE to prepare a random energy eigenstate $|\psi_i\rangle$.

$$|\psi\rangle = \sum_{i} c_{i} |\psi_{i}\rangle \longrightarrow \sum_{i} c_{i} |\psi_{i}\rangle \otimes |E_{i}\rangle \longrightarrow |\psi_{i}\rangle$$

2 Apply a random local unitary transformation

$$|\psi_i\rangle \longmapsto \sum_i c_j |\psi_j\rangle$$
 where $E_j \approx E_i$

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Quantum Metropolis Algorithm (2/3)

3 Attach the energy register $|E_j\rangle$ using QPE and attach the accept/reject qubit.

$$\sum_{j} c_{j} |\psi_{j}\rangle \longmapsto \sum_{j} c_{j} |\psi_{j}\rangle \otimes |E_{j}\rangle \otimes \left(\sqrt{w_{ij}} |1\rangle + \sqrt{1 - w_{ij}} |0\rangle\right)$$

- 4 Measure the accept/reject qubit
 - Accepted! Use QPE to collapse to a new eigenstate $|\psi_j\rangle$ and learn its energy.
 - Rejected! Continue to step 5.

... the implementation

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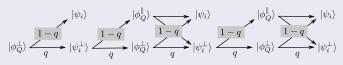
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Quantum Metropolis Algorithm (3/3)

5 Being in state $|\phi_Q^{\perp}\rangle$ we iterate the P Q^{\parallel} measurements in a two dimensional subspace.

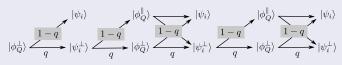


- $|\psi_i\rangle$ hit! Rejection successful, return to step 2.
- $|\psi_i\rangle$ missed! Rejection failed, stop the algorithm.

... the implementation

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